



Utility maximization using Enhanced QoS Parameters on Cloud Storage

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ABSTRACT: Cloud storage(CS) is a model of data storage in which the digital data is stored in logical pools, the physical storage spans multiple servers (and often locations).cloud storage faces conflict in packet delivery rate and delay in transmission of data from user to the cloud storage server. in Cross layer communication. In this paper a Sectorized-Antenna (SA) based protocol is proposed to address the packet delivery rate. The mathematical and experimental results shows proposed scheme is the suggestive alternate that increases in packet delivery rate with reduced average delay that shows the proposed protocol can dealing with QoS requirements.

KEYWORDS: Cloud Computing, QoS parameters, energy efficient, Packet delivery rate and Multipath routing.

I. INTRODUCTION

Cloud storage (CS) is basically attributed of harvesting information from the physical environment, performing trouble-free processing on the extracted data and transmitting it to remote locations. There are many communication protocols designed for data routing in Cloud storage. However, to achieve throughput and utility maximization which complies with the communication capacities, the systems need a central computational point with the knowledge of the actual network structure and parameters. The communication capacity (i.e. throughput) is associated with the links of the graph. TDMA (Time Division Multiple Access) assigns the communication capacities and make definite collision-free communication. The communication in the network is described by communication demands. Each communication demand is given by source nodes; sink nodes and data quantity to be relocated. Consequently data transfer from numerous source nodes to one sink node are describe as one communication demand (multisource). In a parallel way, the replica allows to explain a problem with numerous sink nodes (multi-sink).The problem as a minimum-cost multi product network flow problem, where each commodity represents one communication demand in the Cloud Computing. The majority of the communication protocols are individually urbanized and optimized for different networking layers, i.e., transport, network, medium access control (MAC), and physical layers. While these protocols attain very high performance in terms of the metrics related to each of these entity layers, they are not equally designed and optimized to make best use of the overall network performance while minimizing the energy spending. Considering the scarce energy and processing possessions of CS, cooperative design of networking layers, i.e., cross-layer design stands as the most promising substitute that has gained interest freshly.Accordingly, a cross-layer protocol (XLP) is developed to attain competent and dependable event communication in cloud with minimum energy overheads. In a cross-layer simulation platform, the state-of-the-art layered and cross-layer protocol configurations have been implemented along with XLP to offer an absolute performance evaluation.

II. LITERATURE REVIEW

Cloud storage (CS) have emerged, different optimizations to overcome their constraints in which it has not achieved the severe QoS parameters which require to be esteemed depending on the running application. However, as it is counterproductive for the cloud lifetime, an extra investigation needs to be established out to confirm the performance of the used multimedia coder on the sensors energy consumption. Normally, the work that considers these QoS requirements and coders' only judge packet error rates at the physical level [6]. Energy optimal routing in the network without any central node or knowledge of the whole network structure uses a partial knowledge about the network Furthermore, the asynchronous/synchronous algorithm is not able to adapt the routing subject when the network changes [9].



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Sectorized-Antenna Medium Access Control (SA), an incorporated cross-layer protocol that offer the communication mechanisms for sensor network to entirely utilize sectorized antennas. It obtain a high packet delivery ratio by minimizing channel contention and packet collisions in the common wireless communication medium but not effective in identifying the time slot assignment, time slot duration, and neighbor discovery parameters selection [1]. A distributed, multi constraint, cross-layer QoS routing algorithm for wireless mesh networks that concurrent gratify multiple QoS requirements. The algorithm profitably warranties various QoS requirements and achieves higher network throughput when compared with other standard techniques. A novel integrated QoS performance metric has various QoS constraints are unambiguously considered where the cross-layer QoS routing algorithm cooperate with distributed opportunistic scheduler optimize the network performance [12]. Highly-Functional Cross-Layer Optimized Interfaces will vigorously hold up high bandwidth applications at low cost and with extremely high energy efficiency. The cross-layer nodes permit for quality-ofservice aware packet routing and security, and leverage real-time optical performance monitoring modules to make easy physical-layer aware switching [13]. Cross layer protocol (XLP) accomplish congestion control, routing, and average access control in a cross-layer fashion. The plan standard of XLP is based on the cross-layer concept of initiative determination, which enables receiver-based contention, initiative-based forwarding, local congestion control, and distributed duty cycle operation to recognize efficient and reliable communication in cloud but networking functionalities such as adaptive modulation, error control, and topology control in a cross-layer fashion to expand a combined cross-layer communication module is not achieved [2]. Cross layer optimization algorithm, which includes routing based on neighbor detection and dual congestion control for improving QoS [11] helps in minimize the energy required in broadcast of video packets. Joint optimal design of congestion control and wireless MAC-layer scheduling uses an article generation approach with imperfect scheduling get better the optimal time-share values. Imperfect scheduling attains an approximation ratio that converges to a sub-optimum of the generally problem with the same approximation ratio [15]. Adaptive and distributed completion of policies arising is unspecified that only eagerly available measurements, such as the received data, are available at the receiver in order to play the considered games [5]. Both single-cell and multi-cell networks are considered of non cooperative games for power allocation, spreading code allocation, and choice of the uplink (linear) receiver but does not exist in asynchronous networks operating on multipath channels, a deeper investigation of the resource allocation problem in the multi cell scenario. Adaptive CSMA scheduling algorithm achieves the maximal throughput distributed under some assumptions and is relevant to a very general interference model but end to- end flow control fails in achieving the optimal utility and fairness of competing flows [4]. Though, identify the precise conditions on the step sizes to ensure stability is complicated because they may depend on network size, network topology, and arrival rates. Routing and Spectrum Allocation algorithm through local control actions maximize the network throughput by performing joint routing, active spectrum allocation, scheduling, and broadcast power control [3] but derives only the conjectural lower bound.

III. METHODOLOGY

The survey of research in throughput maximization through cross layer communication is performed. A number of open researches confront addressing which bring considerable benefits to the network users.

1.1 SA Communication protocol for sensor network

Sectorized Antenna Based Medium Access Control (SA) protocol particularly considered with sectorized antenna system in cloud. SA provides the fundamental communication which functionality confines the potential advantages of sectorized antennas [1]. Particularly, SA is designed to attain the subsequent objectives. Despite the possible benefits of sectorized antennas exploiting these benefits is demanding since a communication pair prepared with sectorized antennas must choose the suitable sectors to point to each other during their communication session. Otherwise, the receiving node cannot obtain the message due to weak response power. As a result, insignificant functions like neighbor discovery become extremely challenging. To address these challenges, additional functions should be performed before data forwarding can begin.

SA utilizes a token-passing approach that serializes the neighbor discovery process along with all nodes. When all nodes discover their neighbors, the absolute neighborhood information is gathered at the sink node. The similar overheads are predictable for roughly all protocols that require at least local neighbor discovery. Moreover, these overheads are one-time overheads. Subsequent discoveries and updates are again of similar overheads of other



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protocols. To dependably share out the computed time schedules to all sensor nodes in the network, SA utilizes the token-passing approach again.

During this process, from above equation(1) initial time synchronization is established among the nodes. These operations are achieved once for the whole network at the opening of the network's lifetime, and thus, their overhead is rewarded by the network's long lifetime. During link or node failures, schedule computation and distribution are performed partially on the affected area of the network, thus, minimizing the overhead of such protection events.

1.2 XLP for Efficient Communication in cloud

A cross layer protocol (XLP) achieves jamming control, routing, and medium access control in a cross-layer manner. The design principle of XLP is a combined cross-layering such that both the information and the functionalities of three primary communications patterns (medium access, routing, and congestion control) are considered in a single protocol operation. The XLP enables receiver-based contention, initiative-based forwarding, local congestion control, and distributed duty cycle operation to understand competent and dependable communication in cloud. In CS, the main goal of a communication group is to productively transport event information by constructing (possibly) multihop paths to the sink. Then, cross-layer initiative determination constitutes the core of the XLP and absolutely incorporates the intrinsic communication functionalities required for successful communication in cloud. A new hop-by-hop local cross-layer congestion is included with XLP based on the buffer residence analysis. The objective of this component is to perform hop-by-hop congestion control by exploiting the local information in the receiver-contention and avoid the need for end-to-end congestion control. It also exploits the restricted reliability measures taken by the channel access functionality and therefore, does not require conventional end-to-end reliability mechanisms. The ultimate goal in the cross-layer design techniques is to develop a particular communication module that is accountable for the functionalities of each networking layer. The initiative determination concept developed in this work is the first step in this approach to restore the entire customary layered protocol architecture that has been used so far in CS so that both the information and the functionalities of customary communication layers are blended in a single module. The cloud classification technique and show how it leads to the formulation of a QoS programming

Finally, XLP is the primary protocol that incorporates functionalities of all layers from physical to transport layer into a cross-layer protocol. A cross-layer analytical construction is developed to examine the performance of the XLP. Furthermore, in a cross layer simulation stage, the state-of-the-art layered and cross-layer protocols have been implemented all along with XLP for performance evaluations.

1.3 Cross-layer Routing and Dynamic Spectrum Allocation

A distributed and restricted Routing and Spectrum Allocation (ROSA) algorithm for joint dynamic routing and spectrum allocation for multi-hop cognitive radio networks addresses routing and spectrum assignment with power control under the so-called physical interference model, which computes the interference among secondary users using a SINR-based model. The ROSA algorithm considers and leverages the unique characteristics of cognitive radio including the accessibility of spectrum holes at an exacting geographic position and their possible variability with time. In the ROSA algorithm each cognitive radio makes real-time decisions on spectrum and power distribution based on nearby composed information. Nodes change their transmission power to exploit the link capacity on the selected spectrum portion. Introduce a concept of spectrum hole that considers interference from bordering resulting as well as primary users, and leverage it to optimize resource utilization at a low computational cost. ROSA calculates the next hop opportunistically depending on queuing and spectrum dynamics, according to the range usefulness function. $sign(x)$ between the large margin and small margin.

A practical implementation of the ROSA algorithm relies on a dual radio with a frequent direct channel and a frequency-agile data channel. ROSA algorithm maximizes the network throughput by performing joint routing, active spectrum allocation, expansion, and transmit power control. Specifically, the ROSA algorithm dynamically assigns spectrum resources to make the most of the capacity of links without generating dangerous interference to other users while guaranteeing bounded bit error rate (BER) for the receiver. In addition, the algorithm aims at maximizing the subjective sum of differential backlogs to stabilize the system by giving priority to higher capacity links with high differential backlog. The proposed algorithm is distributed, computationally experienced, and with bounded BER guarantees.



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1.4 Distributed CSMA Algorithm in Wireless Network

A distributed adaptive CSMA (Carrier Sensing Multiple Access) algorithm for a universal interference model is stimulated by CSMA but applied to more common resource sharing problems (i.e., not limited to wireless networks). CSMA algorithm ignores the packet collisions and achieve maximal throughput, if the adaptation is slow enough [4]. The adaptive CSMA algorithm is a modular MAC-layer protocol that works with other algorithms in transport layer and network layer. Adaptive CSMA scheduling algorithm achieves the maximal throughput distributed under several assumptions. Major compensation of the algorithm applied to a very universal interference model; it is simple, distributed and asynchronous. Furthermore, merge the algorithm with end to end flow control to accomplish the optimal utility and fairness of competing flows. The existing proof of the throughput optimality is based on the stationary distribution of the CSMA Markov chain. This is certainly sufficient, but may not be necessary according to our simulations. However, identifying the exact conditions on the step sizes to ensure stability is difficult.

The potential is often represented by a CSMA Markov chain combination of a set of features, i.e., . Hence the distribution can be written as: $\log P(\mathbf{v}) = \sum_{c \in C} \sum_{G \in \mathcal{G}} w_{fi}(\mathbf{v}_c) - \log Z = \mathbf{w} \mathbf{f} \mathbf{v} \cdot (\cdot) - \log Z$

$$c \in C, G \in \mathcal{G} \quad (7)$$

Let X be a set of conditional (or observed) random variables and Y be a set of target (or label) random variables. A conditional Markov networks (G, Φ) is a Markov network which defines the distribution

$$P(\mathbf{y} | \mathbf{x}) = \frac{Z(\mathbf{x})}{Z(\mathbf{x}, \mathbf{y})} \prod_{c \in C} \phi_c(\mathbf{x}_c, \mathbf{y}_c) \quad (8)$$

Where the partition function is dependent on x, i.e., $Z(\mathbf{x}) = \sum_{\mathbf{y}} \prod_{c \in C} \phi_c(\mathbf{x}_c, \mathbf{y}_c)$

IV. SIMULATION

Performance Analysis of Cross Layer Communication in Cloud Computing uses the ns-2 network simulator. In simulation, set up 'n' nodes consistently at randomly surrounded by 1000×1000 squares, with n unpredictable among 100 and 1000 formative the mobile sensor node group patterns. In particular, to accurately estimate the production of the structure in which each node progress to a randomly selected position with a randomly selected velocity amongst a predefined minimum and maximum speed. The affecting mobile sensor networks continue there for a predefined pause time. After the pause time, it then randomly chooses and moves to another location. This random progression is constant during the simulation period. All simulations were performed for 750 simulation seconds, fixed a pause time of 30 simulation seconds and a minimum moving speed of 1.5 m/s of each node.

In the Random Way Point (RWM) model, each node shift to an erratically chosen location with a randomly selected speed between a predefined smallest amount and highest speed. It assumes the normal unit disc bidirectional communication replica and adjust the message range, so that each node will have roughly 40 neighbors on average. RWM use and standard of the total number of mail sent or received per node as calculated of the communication requirements, and measure resiliency by counting the number of times must run the protocol in order to detect a single node replication. The performance analysis of cross layer communication in Cloud Computing is measured in terms of the Throughput, Average Delay and Delivery Rate

V. RESULTS AND DISCUSSION

In order to compare the throughput efficiency of the algorithms used several metrics to evaluate their performance. The first metric is the total throughput obtained while performing the communication in the Cloud Computing. The second performance metric is called the average delay, which is defined as the percentage of time it takes to receive the bit of data from one node to another. It is typically considered in multiples or fractions of seconds. This metric shows the level by which the QoS requirements negotiated between the network users are violated due to the energy-aware resource management. The third metric is the delivery rate which is the average rate at which output is delivered and is expressed as the ratio of output produced to the elapsed time.

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Table 1 Throughput

No. of active sessions (time)	Throughput (Kbps)		
	Distributed CSMA Algorithm	SA Protocol	XLP in cloud
5	1905	1847	1512
10	2748	2975	2345
15	2903	3027	2498
20	3027	2992	2416
25	3128	2691	2128
30	3317	2415	2077
35	3034	2858	2215

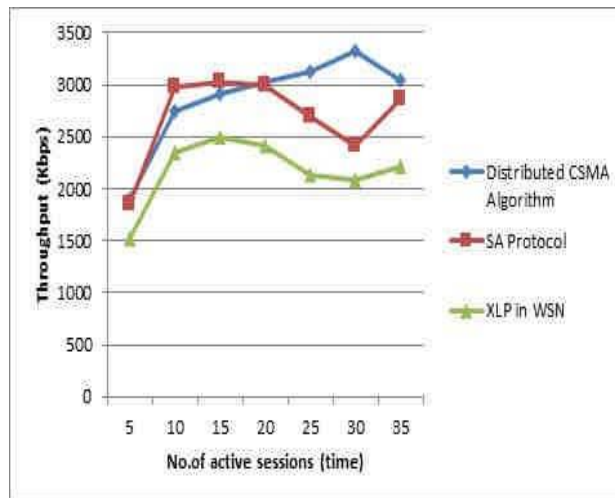


Figure 1 No. of active sessions Vs. Throughput

The above fig 1 describes the throughput ratio in terms of Kilo bits per Second (Kbps) based on the active sessions. The throughput ratio is 10 – 20 % high in Distributed CSMA Algorithm when compare with the SA Protocol and XLP in cloud. The active session for experiment varies from the 5, 10, 15...upto 35. Form the figure it illustrates that the as the session time increases, the throughput increases Distributed CSMA Algorithm, AMAC Protocol and XLP in cloud.



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Table 2: Average Delay

Network Size (nodes)	Average Delay (sec)		
	SA Protocol	XLP in cloud	ROSA Algorithm
18	2	5	10
25	10	14	18
32	12	16	21
40	15	20	23
45	18	22	26
55	22	26	30
60	30	34	39

Table 3: Delivery Rate

No. of Packets	Delivery Rate (data units/ms)	
	ROSA Algorithm	SA Protocol
25	4.02	6.95
50	5.61	8.54
75	6.3	9.38
100	6.24	9.16
125	7.13	10.04
150	8.06	11.87
175	9.15	12.15

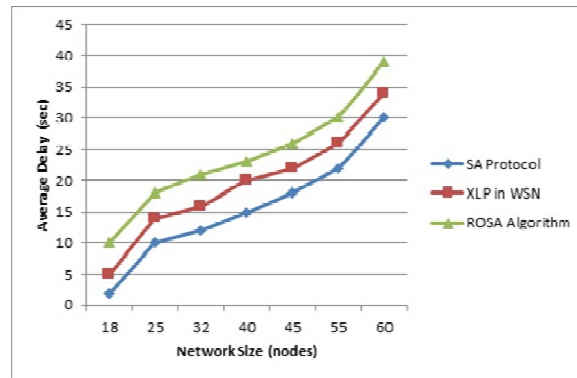


Figure 2: Network Size vs. Average Delay

Fig 2 describes the average delay based on the network size. The average delay is measured in terms of seconds (sec). As the network size increases, delay time is reduced in the SA protocol. It can be seen that the method of this work shows conspicuous advantage. The experiment shows that the method of this work can greatly bring down the time in terms of delay.

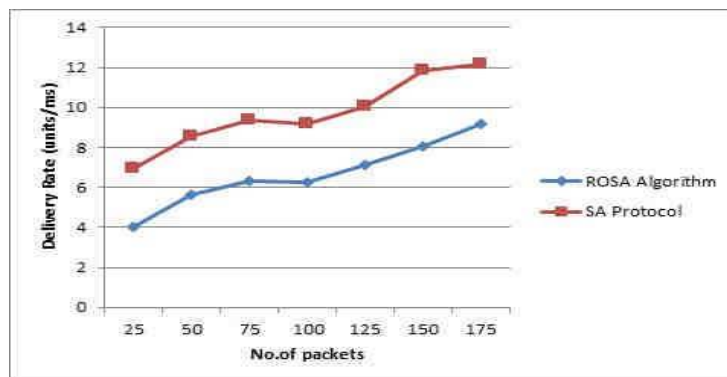


Figure 3 No. of Packets vs. Delivery Rate

The SA protocol delay time is reduced when compared to the XLP in cloud and ROSA Algorithm. SA Protocol is 2 – 5 % lesser delay time taken when compared with the XLP in cloud and 5 -10 % lesser delay time taken when compared with the ROSA Algorithm. SA protocol has lesser interruption while forwarding the packets and performing the communication in sensor network. The results showing the mean delivery rate of SA Protocol and ROSA Algorithm are evaluated. The graph shows that an increase of the packets leads to improved delivery rate. However, the effective delivery rates are achieved with different intervals between the packets.

Compared with SA protocol and ROSA algorithm, it consumes 5 to 10 % higher delivery rate in SA protocol.

VI. CONCLUSION

The efficacy of all the above algorithms is verified by simulations and more widely used in the target tracking applications. In Cloud Computing, an effective communication is achieved using the cross layer communication protocol. The work completely summarizes the cross layer communication strategies and its impacts in Cloud Computing. Some of the strategies discussed above mainly focus on throughput, utility maximization but are lacking in some factors. Hence this survey work will hopefully motivate future researchers to come up with the minimal delay and maximal throughput with utility maximization to strengthen the sensor network paradigm.



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