



Locating Fault with Alarm Correlation Using Link Failure Detection Algorithm in Optical WDM Networks

K.Pandiammal¹, S. Arulselvi*²

Assistant Professor, Dept. of ECE, Jerusalem College of Engineering, Chennai, Tamil Nadu, India¹

Assistant Professor, Dept. of ECE, Bharath University, Chennai, Tamil Nadu, India²

* Corresponding Author

ABSTRACT: For several high speed networks, providing resilience against failures is an essential requirement. The main feature for designing next generation optical networks is protecting and restoring high capacity WDM networks from the failures. Quick detection, identification and restoration make networks more strong and consistent even though the failures cannot be avoided. Hence, it is necessary to develop fast, efficient and dependable fault localization or detection mechanisms. In this paper we propose a new fault localization algorithm for WDM networks which can identify the location of a failure on a failed lightpath. Our algorithm detects the failed connection and then attempts to reroute data stream through an alternate path. In addition to this, we develop an algorithm to analyze the information of the alarms generated by the components of an optical network, in the presence of a fault. It uses the alarm correlation in order to reduce the list of suspected components shown to the network operators. By our simulation results, we show that our proposed algorithms achieve less blocking probability and delay while getting higher throughput.

KEYWORDS: Alarm Correlation, blocking probability, delay, Fault Localization, WDM Networks.

I. INTRODUCTION

A. Wavelength-Division-Multiplexing (WDM) Networks

WAVELENGTH division multiplexing (WDM) is the concurrent transmission of multiple streams of data with the help of the special properties of fiber optics [1]. The ability of transferring huge amount of data at high speeds by the users over large distances is offered by the WDM networks. WDM is considered as the most talented technology for the backbone of future next generation internet [2]. In WDM all optical networks, data is routed through optical channels called light paths. Without the wavelength conversion ability, the light path establishment requires same wavelength to be utilized along the entire route of the light path which is commonly referred to the wavelength continuity constraint. WDM enables the employment of a substantial portion of the available fiber bandwidth by allowing many independent signals with different wavelength to be transmitted simultaneously on one fiber [3]. The routing and detection of these signals can be obtained independently since the wavelength estimates the communication path by acting as the signature address of the origin, destination or routing. Hence, the wavelength selective components are necessary for allowing the transmission, recovery or routing of specific wavelengths.

B. Fault Detection and Localization

At different protocol layers, link failures can be detected. Generally, the detection time is much longer for an upper layer protocol when compared with the optical/physical layer scheme [19]. We focus on optical layer monitoring schemes where a link failure can be detected by a special device called monitor for reducing the detection time. A channel based monitoring scheme requires a large number of monitors because it requires one monitor for each wavelength channel. Though link based monitoring scheme is more scalable, still it requires one monitor per link [4]. When compared with the fault localization, fault detection is easier and faster. Fault localization is the process of finding a minimum set of potential failed network resources based on the alarms generated in the fault detection phase.



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Fault localization in general network has been studied exclusively for many years in various areas and thus it is not a new problem. It has been studied in the areas like power distribution systems, electrical circuits, industrial control systems, and in communication networks. On the other hand, due to the lack of electrical terminations or the excessive cost and the difficulty in implementation, the existing fault localization schemes for traditional networks cannot be applied to the WDM networks directly [5]. In all optical WDM networks, the network edge routers may be able to detect the existence of a fault whenever a link is damaged or a channel is disconnected. But it is not possible to indicate the exact location of the fault. At this time, no advanced optical technique is introduced.

C. Fault Tolerant in WDM Networks

Failures in WDM networks may seriously damage the end user applications because each lightpath can carry a large amount of traffic. Failures in all optical WDM networks can be classified into two categories according to the scale of their effect [6]. The first one is the wavelength level failure which affects the quality of transmission of each individual lightpath. The second one is the fiber level failure which affects all the lightpath on an individual fiber. Failures can lead to a severe data loss because each lightpath is estimated to operate at a rate of several gigabytes per second [7].

The capability of the network to endure the failures is known as fault tolerance. Due to the node failure or link failure, failure arises. When a link tends to fail, then all its constituent fibers also fails. Each and every connection which utilizes these fibers is rerouted and a wavelength will be assigned. Primary path is a light path which carries traffic during normal operation [18]. In case of a failure, the traffic is rerouted. The large amount of traffic on these networks against the traditional copper links makes the fault tolerance as a major issue.

In this paper we propose a new fault localization scheme, which functions in the WDM layer. In this work, we propose a new algorithm which can identify the location of a failure on a failed lightpath. In addition to this, we develop an algorithm to analyze the information of the alarms emitted by the components of an optical network, in the presence of a fault [8-9].

II. PROPOSED FAULT LOCALIZATION AND ALARM CORRELATION

A. Detection and Localization Algorithm

In this section, we propose a Link Failure Detection Algorithm which can identify the location of a failure on a lightpath. Our proposed algorithm detects the failed connections and tries to reroute data stream through an alternate path. Because of the short interval time of establishing a lightpath, we assume that a failure occurs during data transfer. This shows that the destination node knows the source node and the setup route before an interruption occurs [16-17]. When the destination does not receive expected data stream for the given time interval, then the connection gets interrupted. Immediately, our algorithm will be activated by the destination. The operation of our algorithm is given below.

A Link Failure Alarm (LFA) signal is disseminated toward the upstream neighbor backwardly by the destination. The alarming signal is a small control packet, which notifies the upstream node of data disruption. The alarming signal traverses through the secured control network (OSCs) [10]. If the recipient of LFA has not received the data before, it passes the LFA to its upstream neighbor in the path. Otherwise, if the recipient of LFA is also a recipient of data, a reply signal REP is sent to the downstream sender. Once the REP is received, the node will activate a restoration protocol to reroute the affected traffic through an alternate by link restoration. Then the location of the failed connection is identified and it is transmitted to the entire nodes of the network. This is essential in order to maintain correct routing tables and also to prevent blocking of forthcoming calls by the failed connection. Moreover, by notifying the location of the failure, this could accelerate the restoration of longer disconnected light paths, even before they activate any restoration/localization process [11].

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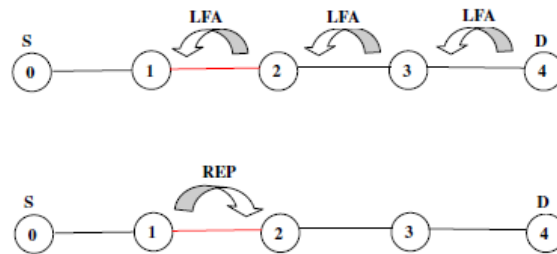


Fig1 Functions of Link Failure Detection Algorithm

1) Algorithm

Let n_1, n_2, \dots, n_k be the set of nodes along the route from the source n_1 to the destination n_k .

1. The source n_1 send data packets in fixed time intervals δ . Source sends packets p_1, p_2, \dots, p_k at time $t_1, t_1 + \delta, \dots, t + k\delta$.
2. The destination n_k checks the packet sequence numbers s_1, s_2, \dots, s_k of the packets.
3. When a $k/2$ number of consecutive packets are missed within a given time threshold t , the destination detects a fault on the light path.
4. If n_k not received the data within the time Interval t , then
 - 4.1 n_k raises LFA and transmit backwards to n_j , where $j = k-1$
 - 4.2 If n_j received the data before t , then
 - 4.2.1 It send a REP packet to n_i , where $i = j+1$.
 - 4.2.2 The n_i starts a recovery scheme for restoration
 - 4.2.3 n_j information is flooded through the network.
 - 4.3 Else
 - 4.3.1 $j = j-1$
 - 4.3.2 Repeat the step 1.2
 - 4.4 End If
5. End If

Fig. 1 shows the functions of our link failure detection algorithm. As shown in the figure, the alarming signal Link Failure Alarm (LFA) is sent back towards the source till it reaches the first recipient of data and it returns the REPLY (REP) signal is returned to the sender. In response, the RFTR scheme [14] is activated to avoid the failed link and restore the connection through a backup path.

B. Alarm Correlation

The components able to send alarm can be divided in the following groups:

- A1 – When the damaged component send alarm;
- A2 – When the component alarms informing that other component is not working correctly.

1) Alarm Correlation Algorithm

When a fault occurs, in order to identify which network component is damaged and which node it belongs each network component has a unique identification. In the network model used here, this identification is composed by a string of four fields

(f_1, f_2, f_3, f_4) having the following meaning for a local node:

- f_1 : It can assume the following values: 0 – non alarming component; 1 – Self alarmed; 2 – Out-alarmed; 3 – Failure masking.
- f_2 : It indicates the node number.
- f_3 : It is always 0 for a local node.



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f_4 : It identifies the position of the component inside the node. The value of this field varies according to the component: LAP = 0 (Local Access Port); ADF = 1 or 2 ; RX = 1 or 2 ; 3R amplifier = 4 or 5 ; TX = 5 or 6 ; PS = 3 .

Algorithm

At the physical route domain, all network components that belong to any channel are numbered as C_1, C_2, \dots and associated to each one of the alarm components of the respective alarms that will be sent to the network management server (NMS) if they fail.

Let S_1 denote the suspected list of components.

1. Server receives alarm from the component C_i .
2. If $a = A_1$ then
 - 2.1 Add $C_i(a)$ to S_1
3. Else If $a = A_2$ then
 - 3.1 For each channel CH_i
 - 3.1.1 Add $C_1(a)$ to S_1 .
- 3.2 For each component $C_k(a)$, $k \neq 1$
 - 3.2.1 If $AD_k = PD(i)$, then where AD_k is alarm domain and PD is the route domain,
 - 3.2.1.1 Add $C_k(a)$ to S_1 .
 - 3.2.2 Else
 - 3.2.2.1 Drop the alarm
 - 3.2.3 End If
- 3.3 End For
4. End If

II. SIMULATION RESULTS

A. Simulation Model and Parameters

In this section, we examine the performance of our Fault Localization and Alarm Correlation algorithm (FLAC) with an extensive simulation study based upon the ns-2 network simulator [15]. We use the Optical WDM network simulator (OWNs) patch in ns2, to simulate an 8-Node topology (Fig.3). Various simulation parameters are given in table I.

TABLE I
SIMULATION PARAMETERS

Topology	Mesh
Total no. of nodes	8
Link Wavelength Number	8
Link Delay	10ms
Wavelength Conversion Factor	1
Wavelength Conversion Distance	8
Wavelength Conversion Time	0.024
Link Utilization sample Interval	0.5
Traffic Arrival Rate	0.5
Traffic Holding Time	0.2
Packet Size	200
No. of Session-traffics	4
Max Requests Number	50
No. of Link Failures	1,2,3 and 4

In our experiments, we use a dynamic traffic model in which connection requests arrive at the network according to an exponential process with an arrival rate r (call/seconds). The session holding time is exponentially distributed with mean holding time s (seconds).

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In our experiments, we will consider multiple link failures. We vary the number of link failures as 1, 2, 3 and 4. The connection requests are distributed randomly on all the network nodes. In all the experiments, we compare the results of FLAC with End-to-End [5] scheme.

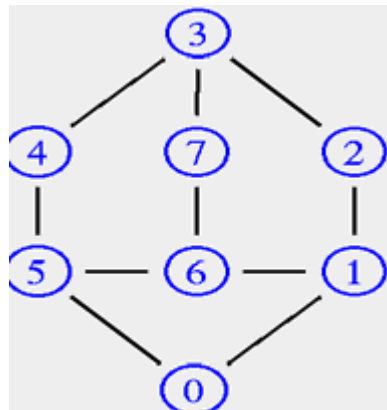


Fig. 2 8-Node Topology

B. Performance Metrics

In our experiment, we measure the blocking probability, end-to-end delay and throughput.

Blocking Probability: It is the ratio of number of rejected requests to the total number of requests sent.

Average end-to-end delay: The end-to-end-delay is averaged over all surviving data packets from the sources to the destinations.

Throughput: It is the number of packets received successfully.

Channel Utilization: It is the ratio of bandwidth received into total available bandwidth for a traffic flow.

C. Results

Our simulation results for the number of failures are presented in Table II.

TABLE II
SIMULATION RESULTS

No of failures	Blocking Probability		Delay		Throughput	
	FLA C	End-To-End	FLA C	End-To-End	FLA C	End-To-End
1	0.5432	0.68667	1.0966	2.32457	732	427
2	0.6178	0.71666	1.7354	3.09908	642	281
3	0.6805	0.78574	2.2553	4.72548	564	152
4	0.7416	0.82366	4.4825	6.57268	405	104

Our simulation results for the various rates are presented in Table III.

TABLE III
SIMULATION RESULTS

Rate (Mb)	Blocking Probability	Delay	Utilization
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	FLAC	End-To-End	FLAC	End-To-End	FLAC	End-To-End
2	0.01	0.15	200.8101	258.462486	0.034098	0.015658
4	0.122075	0.2093846	165.6041	243.106743	0.030568	0.01366
6	0.304862	0.4923077	101.2084	230.450229	0.024459	0.006868
8	0.507692	0.6967554	59.88886	152.666943	0.017623	0.002062

A. Based on Failures

In the initial experiment, we vary the number of failures as 1, 2, 3 & 4 and measure the blocking probability, end-to-end delay and Throughput [12-13].

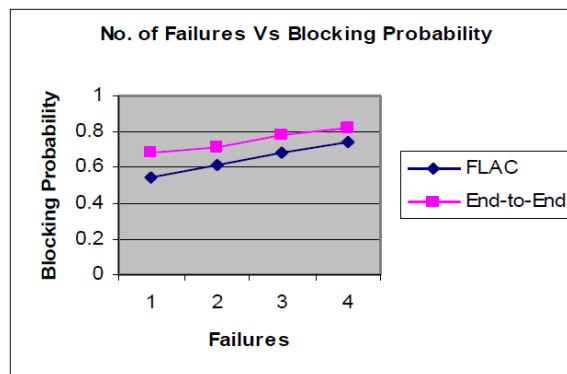


Fig. 3 No. of Failures Vs Blocking Probability

Fig.3 shows the blocking probability obtained with our FLAC algorithm compared with End-to-End scheme. It shows that the blocking probability is significantly less than the End to- End, when number of failures increases.

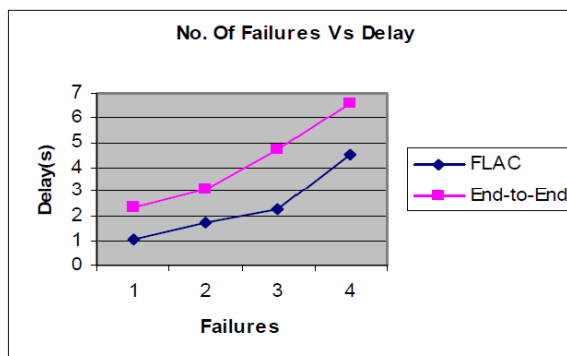


Fig. 4 No. of Failures Vs Delay

Fig.4 shows the end-to-end delay occurred for the number of failures. It shows that the delay of FLAC is significantly less than the End-to-End scheme.

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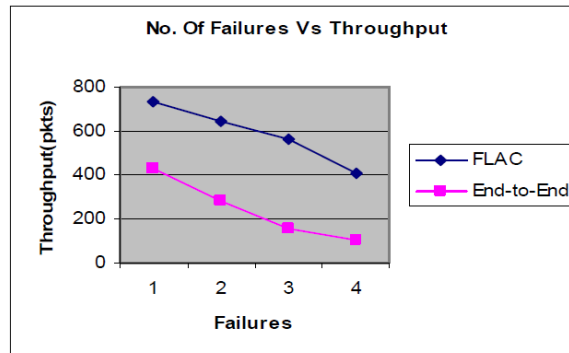


Fig. 5 No. of Failures Vs Throughput

Fig. 5 shows the throughput occurred for various failures. As we can see from the figure, the throughput is more in the case of FLAC when compared to End-to-End scheme.

B. Based on Rate

In the next experiment, we vary the traffic rate as 2Mb, 4Mb....8Mb and measure the blocking probability, end-to-end delay and channel utilization.

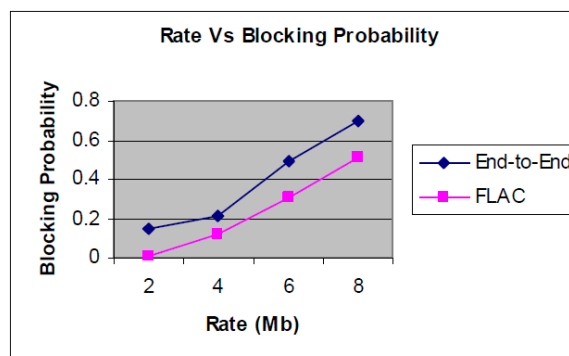


Fig. 6 Rate Vs Blocking Probability

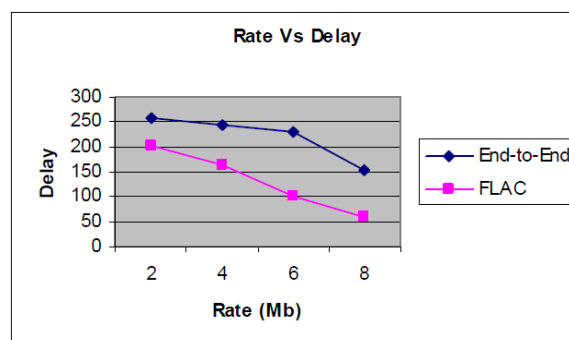


Fig. 7 Rate Vs Delay

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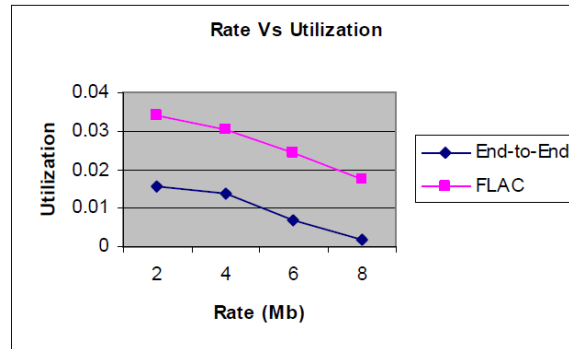


Fig.6 shows the blocking probability obtained with our FLAC algorithm compared with End-to-End scheme. It shows that the blocking probability is significantly less than the DPBR, as rate increases.

Fig.7 shows the end-to-end delay occurred for various rates. It shows that the delay of FLAC is significantly less than the End-to-End scheme.

Fig.8 shows the channel utilization obtained for various rate. It shows that FLAC has better utilization than the End-to-End scheme.

V. CONCLUSION

Quick detection, identification and restoration make networks more strong and consistent even though the failures cannot be avoided. Hence, it is necessary to develop fast, efficient and dependable fault localization or detection mechanisms. In this paper we have proposed a new fault localization algorithm for WDM networks which can identify the location of a failure on a failed lightpath. Our algorithm detects the failed connection and then attempts to reroute data stream through an alternate path. We have assumed that a failure happens during the data transfer mode due to the short interval time of establishing a lightpath. This implies that the destination node is aware of the source node and the setup route before an interruption of service occurs. The algorithm will be activated once a connection is disrupted. This happens when the destination does not receive expected data stream any longer. In addition to this, we have developed an algorithm to analyze the information of the alarms generated by the components of an optical network, in the presence of a fault. The algorithm proposed uses the alarm correlation in order to reduce the list of suspected components shown to the network operators. By our simulation results, we have shown that our proposed algorithms achieve less blocking probability and delay while getting higher throughput.

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