



# **Simplified AODV with Urgent Data Indication (sEAODV<sub>u</sub>): An Energy Efficient Routing Protocol for Zigbee Network**

Sohil A. Dabhi<sup>1</sup>, Manish K. Patel<sup>2</sup>

Assistant professor, Dept. of EC, Dharmsinh Desai University, Nadiad, Gujarat, India<sup>1</sup>

Assistant professor, Dept. of EC, Dharmsinh Desai University, Nadiad, Gujarat, India<sup>2</sup>

**ABSTRACT:** ZigBee is the emerging industrial standard for Ad-hoc wireless networks based on IEEE 802.15.4, having characteristics such as low data rate, low price, and low power consumption. It is widely used in wireless sensor networks for remote monitoring, Home-appliance networks, and industrial automation. It supports tree and mesh routing. A major concern in wireless sensor networks (WSNs) is energy conservation, since battery-powered sensor nodes are expected to operate for a long time (for months or even years). AODV, an on-demand routing protocol is used for ZigBee mesh networks. The high number of control packets, makes AODV unsuitable for certain applications that requires energy efficient communications. A new routing strategy (sEAODV) is proposed here: combination of i) simplified AODV (sAODV): will try to reduce control overhead and hence energy consumption and ii) Energy Aware Mesh Routing (EAODV): to prevent early death of nodes having low energy (less than 10% of its initial energy level) will not forward RREQ messages hence, they will not be a part of communication. The algorithm performance is simulated using NS-2 simulator. Our proposed scheme shows overall performance improvement compared to original AODV in terms of Energy Consumption, Routing Overhead and Number of packet drops. At the end one idea is proposed to indicate emergency/urgent data (fire alarm or sensory data of a patient in ICU) by a flag bit, which forces even the low energy node to deliver the data to the destination node.

**KEYWORDS:** ZigBee, Energy Efficient, Simplified AODV, NS-2 simulator, Wireless Sensor Networks, Routing.

## **I.INTRODUCTION**

The Ad-hoc On-demand Distance Vector (AODV) routing protocol provides a method of routing in mobile ad-hoc networks. Though AODV is simple in comparison to other MANET routing protocols, the specification still contains many sections prone to erroneous programming. AODVjr is a trimmed down AODV, which removes some elements of AODV. It is shown[2] that AODVjr has nearly the same performance as AODV.

**AODVjr[2]** is a Simplified Version of AODV. As AODV was designed to incorporate many features and designed to maximize performance, so it is bit complex on the other hand AODVjr does not require:

1. RERR messages
2. Sequence numbers
3. Precursor lists
4. Gratuitous RREP
5. Hop count
6. Hello messages.

AODVjr algorithm is an on-demand routing algorithm and has many great advantages in energy-saving and network performance. The algorithm only keeps the valid route and does not maintain the routes which couldn't arrive at the routing destination during communications. Removing sequence numbers results in only destination to respond to RREQ. This also eliminates Gratuitous RREP since all routes will be bidirectional. Since the destination will only respond to the first RREQ it receives the "best" (fastest) route is always chosen regardless of the number of hops.



# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

Vol. 3, Issue 10, October 2014

To perform route maintenance route lifetimes are only updated by the reception of packets and not the sending of packets. This requires the destination to occasionally send a packet to the source. If data traffic is unidirectional periodic messages (*connect*) are sent to maintain the route. If data communications are bidirectional, no additional overhead is needed. Using this end-to-end strategy, hello messages, RERR and precursor lists are not needed.

## II. RELATED WORK

To decide on the optimal path in AODVjr-network, link quality is used instead of hop count. This is achieved by means of calculating the link cost  $C(l)$  and path cost  $C(P)$  in accordance with Eq. (1) and Eq. (2)[3], where  $p$  is the probability of successful packet delivery in link  $l$ .

$$C(l) = \min \left[ \left( 7, \text{round} \left( \frac{1}{p^4} \right) \right) \right] \dots (1)$$

$$C(P) = \sum_{i=1}^{L-1} C(l_i) \dots (2)$$

But it is an energy-unaware algorithm which may cause a high energy consumption of nodes leading to the early death of the whole networks.

In order to solve this problem, Ref. [3] presented energy-aware routing. This paper proposes a multi-path energy balance routing algorithm E-AODVjr[3] to reduce energy consumption, improve network energy efficiency as well as avoid network segmentation and node death. When Energy Balance is used to select a path to transmit data, the appropriate nodes with more energy will be selected for it.

Also EA-AODV[4] built optimal path dynamically according to the remaining energy of the node in order to decrease and balance network energy consumption. As the speed of energy consumption of each node is different, all the nodes are divided into three different energy levels which can react to the routing request according to their own energy level in the process of data transferring.

## III. PROBLEMS

### Case [A]

During the route discovery phase in AODV and AODVjr, a node must forward each RREQ packet received by it regardless of its remaining battery/energy level. This is a fatal waste of energy for nodes having lower energy level. As these nodes will go to die soon, just after they become a relay node in forwarding packets from sender to the destination as shown in Figure 1.

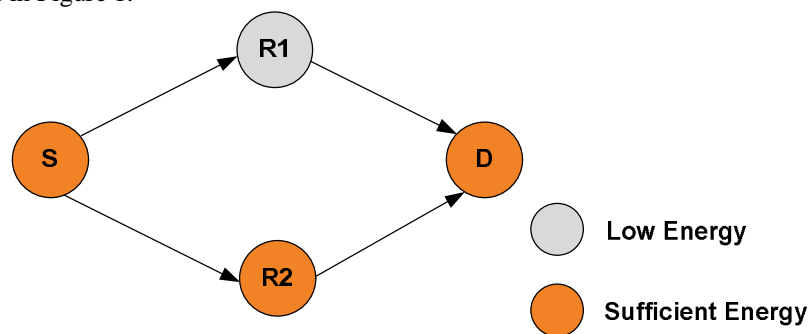


Figure 1: RREQ forwarded by R1 would become useless soon

Node S (sender) starts route discovery to find the optimal route to Node D (destination) as shown in Figure 1. There are two paths available: Path1 (S - R1 - D) and Path2 (S - R2 - D). Node R1 with low energy and Node R2 with sufficient energy broadcasts the RREQ packet received during route discovery phase. However, due to the low path cost of Path

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

Vol. 3, Issue 10, October 2014

1, it turns out to be the optimal route. The problem is: As the energy level of Node R1 is low, it drains all the energy very soon. And that results in a route failure and necessity of finding new route will arise. To solve the problem mentioned above, we propose an improved AODV algorithm which states that if a node is having low remaining energy (less than 10%) then it should not participate in forwarding in RREQ, hence it will not be a relay node for forwarding packet/s.

### Case [B]

Node S starts route discovery to find the optimal route to Node D. There is only one path available: Path1: S - R1 - D as shown in Figure 2.

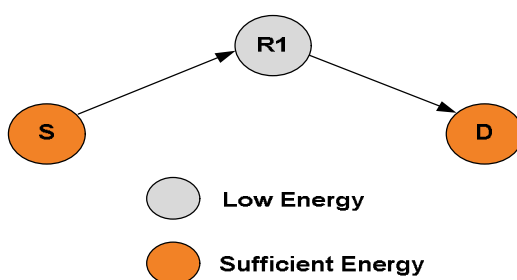


Figure 2: Only one route is available via low energy node

Node R1 with low energy (less than 10%) will not forward the RREQ and hence it will not deliver a packet to D. As a consequence, data cannot be sent from S to D. To avoid this problem in case of emergency or urgent data (fire alarm or sensory data of a patient in ICU) the sending node will use one *flag bit* in control frame of network layer. This bit location is shown in Figure 5.6. If the value of this bit is set to 1 then any node will try its best to send the data to the destination node in spite of having low energy level. But if this bit value is 0 then energy aware mechanism as mentioned above in Case [A] will be applied and the node will not be a relay node for forwarding packet/s.

## IV. PROPOSED ALGORITHM

AODV has some drawbacks and some modifications are already proposed by different authors to improve the performance of the same. Further, we have seen AODVjr, which is actually simplified AODV, also performs almost similar to AODV[2]. Figure 3 shows the flow of proposed modifications. First of all Original AODV protocol is simplified in block(2). This modified code removes the following parameters from the original AODV: -Gratuitous RREP, Hello messages, and Precursor lists. To add the remedy for energy aware routing the protocol is further modified as shown in block(3). Which makes the route finding procedure more efficient. As discussed in Case: [A], if a node is having low remaining energy (less than 10%) participates in forwarding RREQ then it will die out soon. So nodes running on low energy are not allowed to forward RREQ. Which makes them automatically inactive and we can prevent the death of those nodes in near future.

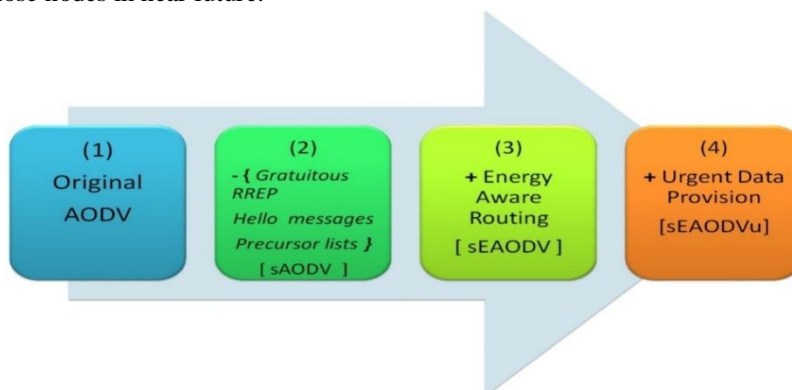


Figure 3: Flow of proposed modification



# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

Vol. 3, Issue 10, October 2014

## V.RESULT AND DISCUSSION

To compare the performance of original AODV with sEAODV, the NS-2 simulator was utilized. The basic AODV code was a version updated by Marina Mahesh. The Simplified AODV-sAODV code was developed for this testing and energy aware routing mechanism is also included, which forms the sEAODV version.

Various fields were used -

- 250 m x 250 m field with 25 nodes,
- 500 m x 500 m field with 50 nodes,
- 700 m x 700 m field with 75 nodes,
- 1000 m x 1000 m field with 100 nodes,
- 1250 m x 1250 m field with 150 nodes,
- 1500 m x 1500 m field with 200 nodes.

The nodes placed randomly within the simulation area. For each result, simulation is performed 50 times and the average is taken for the same. The data rate is 250 kbps, packet size is 50 bytes and simulation time is 200 sec.

Results: Figure 6, 7, 8, 9.

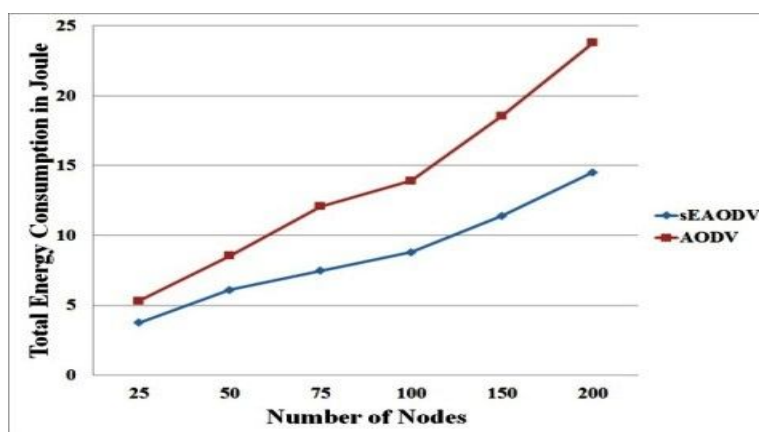


Figure 6: Total Energy Consumption in Joule

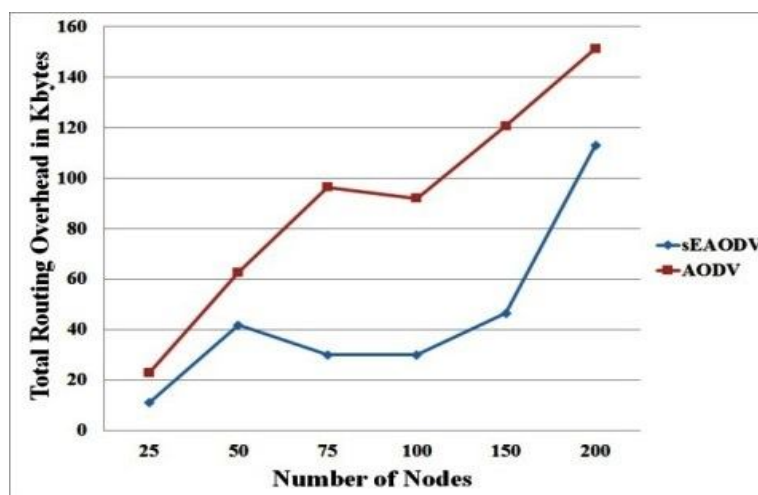


Figure 7: Total Routing Overhead in Kbytes

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

Vol. 3, Issue 10, October 2014

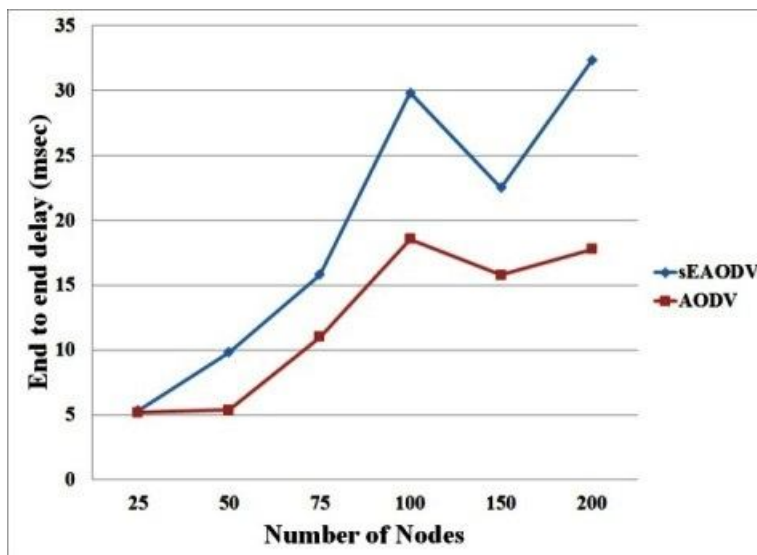


Figure 8: End to End Delay in msec

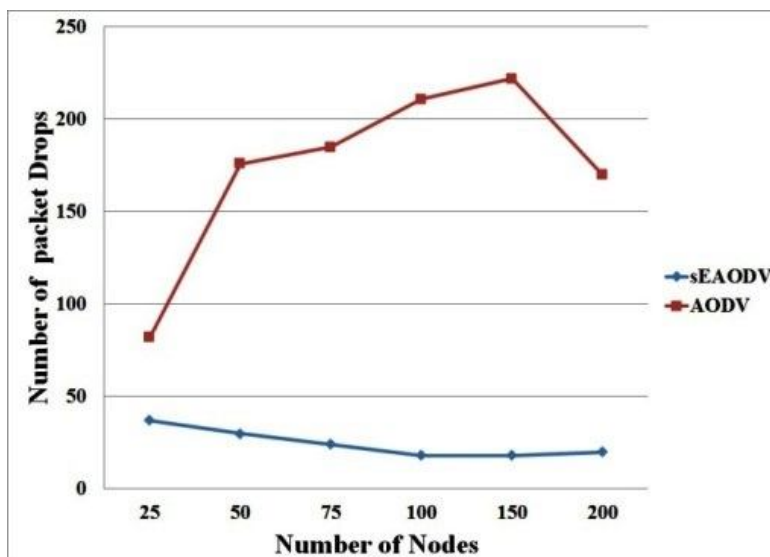


Figure 9: Number of Packet Drops

The set of simulations compares AODV and sEAODV with varying number of nodes. The high routing overhead in AODV is due to the broadcast of hello messages by every node. Hello broadcast messages must be processed by each receiving node. Figure 6 shows the Energy consumption in Joule by entire network. sAODV outperforms AODV when the network size is large. Figure 9 shows the number of packet drops due to nodes with low remaining energy level acting as relay node for communication link and dies out early, which results in link failure and hence drop of packet. Along with increasing network size congestion will also increase.

## VI. CONCLUSION

As we can see from the above results sEAODV performs better than original AODV. Which shows overall performance improvement compared to original AODV. Total Energy Consumption is reduced up to 35%. The performance of the route discovery has also been improved, not only in Total Routing Overhead (around 51% reduction), but also in Number of packet drops (around 74% reduction). But the drawback is, we are compromising on



ISSN (Print) : 2320 – 3765  
ISSN (Online): 2278 – 8875

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

Vol. 3, Issue 10, October 2014

end to end delay. Which is increased by 58% compared to original AODV. At the end, one idea is proposed to deliver urgent packet/s even by the nodes which are dying (having low battery level). For this, one flag bit of control frame at network layer can be used.

## • Scope of Future Work

Application layer implementation of this urgent flag bit (present in Control Frame at Network Layer) and result analysis. Further the battery level of a node can be divided in more number of levels and accordingly multiple tasks can be assigned to the node based on current remaining battery level while processing RREQ.

## REFERENCES

- [1] Taehong Kim, Daeyoung Kim, Noseong Park, Seong-eun Yoo, Toms Snchez Lpez, "Shortcut Tree Routing in ZigBee Networks", IEEE, Information and Communications University, Electronics and Telecommunications Research Institute, 1-4244-0523-8/07/2007.
- [2] Ian D. Chakeres, Luke Klein-Berndt, "AODVjr, AODV Simplified", Mobile Computing and Communications Review, Volume 6, Number 3.
- [3] Jun Xiao and Xiaojun Liu, "The Research of E-AOMDVjr Routing Algorithm in ZigBee Network", IEEE, Chinese Control and Decision Conference, 978-1-4244-8738-7/11/26.00, pp. 2360-2365, 2011.
- [4] Yigong, P., Yingli, L., Hongcheng, L. & Jinshou, Y., "Method for saving energy in ZigBee network", Wireless Communications, pp. 24-26, September 2009.
- [5] Antonio M. Ortiz, Fernando Royo, Teresa Olivares and Luis OrozcoBarbosa, "Intelligent Route Discovery for ZigBee Mesh Networks", IEEE, 978-1-4577-0351-5/11/ 2011.
- [6] V. RenuKachhoria, Renu Kachhoria, "An Implementation in AODV based on Active Route Timeout between sensor nodes in Wireless Sensor Networks", International Journal on Computer Science and Engineering (IJCSSE), Vol. 3 No. 3 Mar 2011.
- [7] Maged Salah Eldin Solimana, Sherine Mohamed Abd El-kaderb, Hussein Sherif Eissac, Hoda Anis Baraka, "New Adaptive Routing Protocol for MANET", Ubiquitous Computing and Communication Journal, Volume 2 Number 3.
- [8] Teerawat Issariyakul, Ekram Hossain, "Introduction to Network Simulator NS2", Springer Science+Business Media, LLC, 2009 (E-book).
- [9] Shahin Farahani, "ZigBee Wireless Networks and Transceivers", Copyright 2008, Elsevier Ltd. ISBN: 978-0-7506-8393-7 (E-book).