# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering 

(An ISO 3297: 2007 Certified Organization)
Vol. 3, Issue 4, April 2014

# Design and Implementation 16 BIT REVERSIBLE LOGIC ALU with 15-Operations 

Abdul Kareem ${ }^{1}$, Mrs. E Kavitha ${ }^{2}$<br>PG Student [M.Tech in Electronics], Department Electronics \& Communication Engineering, HKBK College of Engineering, Bangalore, India ${ }^{1}$<br>Research scholar, St. Peter's institute of Higher education and research, St. Peter, University Avadi, Chennai, Professor

PG-coordinator, HKBKCE, Bangalore, Karnataka, India ${ }^{2}$


#### Abstract

Reversible logic is one of the emerging technologies having promising applications in quantum computing. This project will deal with the design of a 16 bit reversible Arithmetic Logic Unit (ALU) with 15 operations is presented by making use of Double Peres gate, Fredkin gate, Toffolli gate, DKG gate and NOT gate.

A new VLSI architecture for ALU using reversible logic gates is proposed. ALU is one of the most important components of CPU that can be part of a programmable reversible computing device such as a quantum computer.

A first single bit reversible ALU and second single bit ALU are designed and Then 16 single bit ALU's are cascaded together taking carry out of ALU performing LSB operation as an input to carry in of ALU performing next LSB operation. Design is implemented and verified in Verilog in modelsim Altera 6.6d.


KEYWORDS: Reversible logic, Reversible ALU, MODEL SIM, Moore's law.

## I.INTRODUCTION

It is well known that Moore's law will stop to function sooner and something dramatic will therefore have to happen in microelectronics in near future. With much faster and more complex digital systems being built, power consumption of CMOS circuits has become a major concern.

Landauer [2] proved that power loss is an integral feature of irreversible circuits that have information loss irrespective of the technology the circuit is implemented in.

Also, Bennett [3] showed that in order to keep a circuit from dissipating any power, it had to be composed of reversible gates.

Reversible are circuits (gates) that have the same number of inputs and outputs and there is a one-to-one mapping between vectors of inputs and outputs. Thus the vector of input states can be always uniquely reconstructed from the vector of output states. Because truly low power circuits cannot be built without the concepts of reversible logic, various technologies and circuits for reversible logic are recently being studied [2].

The Arithmetic Logic Unit (ALU) is essentially the heart of a CPU. This allows the computer to add, subtract, and to perform basic logical operations such as AND, OR etc. Since every computer needs to be able to do these simple functions, they are always included in a CPU. An ALU is a combinational logic circuit that can have one or more inputs and only one output. ALU's output is dependent only on inputs applied at that instant as a function of time, and not on past conditions. A simple ALU in its basic form consists of two inputs for the operands, one input for selecting the desired operation and one output for the result. The complexity of ALU may vary from processor to processor. In [1], a reversible ALU for one binary arithmetic and three logical operations is designed in Base paper. In present work 16 bit ALU with 15 operations is design.

## II. RELATED WORK

Reversible Arthematic logic unit [1] with 4-operation AND, OR, X-OR and ADD. 4-bit adder/subtractor [7], Design of a 4-bit 2's Complement Reversible Circuit [9], Design of Control Unit for Low Power ALU with a Barrel Shifter Using Reversible Logic [10], Design of 32 Bit Reversible ALU [11] with 7-operations and Arithmetic \& Logic Unit (ALU), Design using Reversible Control Unit [12] with 9-operations are related and recent work in the field reversible logic circuits.

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering <br> (An ISO 3297: 2007 Certified Organization) 

Vol. 3, I ssue 4, April 2014

## III. REVERSIBLE GATES

Reversible logic gates Used in Design of circuit in this paper are NOT gate [8], Feynman Gate [4], Toffoli gate [5], Fredkin Gate [6], Double Peres gate [1] \& DKG gate [7].
NOT Gate [8] is a simplest Reversible gate and is a 1*1 gate. Not gate is shown in the Figure 1 and its quantum cost is Zero.


FEYNMAN GATE [4] Controlled NOT (CNOT) gate is an example for a $2 * 2$ gate. The Reversible $2 * 2$ gate with Quantum Cost of one. It is shown in the Figure 2


TOFFOLI GATE [5] 2 CNOT gatesis $3 * 3$ Reversible gate with three inputs and three outputs. Its quantum cost is 5 and is shown in Figure 3


Figure-3. Toffolli Gate
FREDKIN GATE [6] is $3 * 3$ gate maps inputs $(A, B, C)$ to outputs ( $\mathrm{P}=\mathrm{A}, \mathrm{Q}=\mathrm{A}^{\prime} \mathrm{B}+\mathrm{AC}, \mathrm{R}=\mathrm{AB}+\mathrm{A}^{\prime} \mathrm{C}$ ) having Quantum cost of 5 and It is shown in the Figure 4


Figure-4. Fredkin Gate
DOUBLE PERES GATE [1] is $4 * 4$ gate with quantum cost of 6 . It is shown in the Figure-5

## International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)
Vol. 3, Issue 4, April 2014


DKG GATE [7] is 4* 4 reversible DKG gate [6] that can work singly as a reversible Full adder and a reversible Full subtractor is shown in Figure 6. It can be verified that input pattern corresponding to a particular output pattern can be uniquely determined. If input $A=0$, the proposed gate works as a reversible Full adder, and if input $A=1$, then it works as a reversible Full subtractor.


Figure-6. DKG Gate

## IV. PROPOSED WORK

## 1-BIT-ALU

The ALU that is proposed is 15 -operations. ThereALU has 2 parts. $1^{\text {st }}$ which has Double Peres Gate [1] as base of the circuit and is selected when select line s3 is zero. The operations performed here are buffer, AND, OR, NAND, NOR, EX-OR, and EX-NOR. $2^{\text {nd }}$ part has DKG Gate as base of the circuit and is selected when select line s3 is one. The operations performed here are add, increment, 2's complement, set, subtract, decrement, not, and clear. The operations selected depending on various select line are shown in the table-1

## International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)
Vol. 3, Issue 4, April 2014

| S3 | S2 | S1 | S0 | Operations |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | AND |
| 0 | 0 | 0 | 1 | NAND |
| 0 | 0 | 1 | 0 | OR |
| 0 | 0 | 1 | 1 | NOR |
| 0 | 1 | 0 | 0 | BUFFER A |
| 0 | 1 | 0 | 1 | EX-OR |
| 0 | 1 | 1 | 0 | BUFFER B |
| 0 | 1 | 1 | 1 | EX-NOR |
| 1 | 0 | 0 | 0 | ADDITION |
| 1 | 0 | 0 | 1 | INCREMENT |
| 1 | 0 | 1 | 0 | $2{ }^{\prime}$ s COMPLEMENT |
| 1 | 0 | 1 | 1 | SET |
| 1 | 1 | 0 | 0 | SUBSTRACTION |
| 1 | 1 | 0 | 1 | DECREMENT |
| 1 | 1 | 1 | 0 | NOT |
| 1 | 1 | 1 | 1 | CLEAR |

Table-1


Figure-7. 2:1 MUX
The 2:1 Multiplexer is designed using Fredkin gate when we make A as select line and ( $\mathrm{B} \& \mathrm{C}$ ) as input. B or C is selected depending on A is 0 or 1 respectively. The block diagram is shown in figure 7.
Design of $1^{\text {st }} 1$-bit ALU with 15 operations is shown in the figure 8.


Figure-8. $1^{\text {st }} 1$-bit ALU
For designing of 16 -bit ALU we cascade 161 -BIT-ALU as shown in the figure-10. We need to note here that only $1^{\text {st }}$ bit has to be added with one for increment and 2's complement or subtracted for decrement. There is slight

## International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering

(An ISO 3297: 2007 Certified Organization)

## Vol. 3, Issue 4, April 2014

change for SET and CLEAR operation also. Hence a different 1-BIT-ALU is designed as shown in figure 9 and cascaded from $2^{\text {nd }}$ bit onwards till 16 -bit. It is from (A1-A15) in the figure 10 . A0 in the figure 10 is circuit shown figure 8 .


Figure-9. $2^{\text {nd }} 1$-BIT-ALU


Figure-10. 16-BIT-ALU

## V. SIMULATION RESULT

The simulation result for code for 16-bit-ALU design written in Verilog and verified in model simaltera 6.6d is shown in figures (11, 12, 13 and 14).


Figure-11. Result for $1^{\text {st }} 4$-operations (AND, NAND, OR and NOR)
The Figure-11 shows the result of first 4 operations (AND, NAND, OR and NOR). These are operations are selected depending on select lines s3-s0 as shown in table-2. The Graph is shown each 4 -operations because it can't be

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering 

## (An ISO 3297: 2007 Certified Organization)

Vol. 3, I ssue 4, April 2014
accommodated in the same screen. In Present work the ALU is of 16 -bit and has 15 -operational features were as in the existing work the simulation result and design is for 8 to 10 features at the max.

| S3 | S2 | S1 | S0 | Operation | Input(A) | Input(B) | OUTPUT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | AND | 1100110011001100 | 1010101010101010 | 1000100010001000 |
| 0 | 0 | 0 | 1 | NAND | 1100110011001100 | 1010101010101010 | 0111011101110111 |
| 0 | 0 | 1 | 0 | OR | 1100110011001100 | 1010101010101010 | 1110111011101110 |
| 0 | 0 | 1 | 1 | NOR | 1100110011001100 | 1010101010101010 | 0001000100010001 |

Table-2 shows the first 4 operations selected depending on select lines. The operation is performed on 2 input's $A \& B$ and output is given out.


Figure-12. Results for $2^{\text {nd }} 4$-operations (Buffer A, EX-OR, Buffer B, EX-NOR)
The Figure-12 shows are result of second 4-operations (Buffer-A, EX-OR, Buffer B, EX-NOR).These are selected depending on select lines s3-s0as shown in the table-3

| S3 | S2 | S1 | S0 | Operation | Input(A) | $\operatorname{Input(B)}$ | OUTPUT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 0 | Buffer-A | 1100110011001100 | 1010101010101010 | 1100110011001100 |
| 0 | 1 | 0 | 1 | EX-OR | 1100110011001100 | 1010101010101010 | 0110011001100110 |
| 0 | 1 | 1 | 0 | Buffer-B | 1100110011001100 | 1010101010101010 | 1010101010101010 |
| 0 | 1 | 1 | 1 | EX-NOR | 1100110011001100 | 1010101010101010 | 1001100110011001 |

Table-3
Table- 3 shows the second 4 operations selected depending on select lines. The operation is performed on 2 input's A \& B and output is given out. The operations performed are Buffer-A, EX-OR, Buffer B, and EX-NOR.


Figure-13. Results for $3^{\text {rd }} 4$-operations (ADD, INCREMENT, 2'S COMPLEMENT, SET)

# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering 

## (An ISO 3297: 2007 Certified Organization)

## Vol. 3, Issue 4, April 2014

The Figure-13 shows are results of Third 4-operations (ADD, INCREMENT, 2'S COMPLEMENT, and SET).These are selected depending on select lines s3-s0 as shown in the table-4

| S3 | S2 | S1 | S0 | Operation | Input(A) | $\operatorname{Input}(B)$ | OUTPUT | Co |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 0 | ADD | 1100110011001100 | 1010101010101010 | 0111011101110110 | 1 |
| 1 | 0 | 0 | 1 | INCREMENT A | 1100110011001100 | 1010101010101010 | 1100110011001101 |  |
| 1 | 0 | 1 | 0 | $2 ’$ S <br> COMPLEMENT | 1100110011001100 | 1010101010101010 | 0011001100110100 |  |
| 1 | 0 | 1 | 1 | SET | 1100110011001100 | 1010101010101010 | 1111111111111111 | 1 |

## Table-4

Table- 4 shows the Third 4 operations selected depending on select lines. The operation is performed on 2 input's A \& B and output is given out here Co is consider for ADD \& SET operations.


Figure-14. Results for $4^{\text {th }} 4$-operations (SUBTRACTION, DECREMENT, NOT, CLEAR)
The Figure-14 shows the result of Fourth 4-operations (SUBTRACTION, DECREMENT, NOT, CLEAR). These are selected depending on select lines s3-s0 as shown in the table-4

| S3 | S2 | S1 | S0 | Operation | Input(A) | $\operatorname{Input(B)}$ | OUTPUT | Co |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 0 | 0 | SUBSTRACTION | 1100110011001100 | 1010101010101010 | 0010001000100010 | 0 |
| 1 | 1 | 0 | 1 | DECREMENT A | 1100110011001100 | 1010101010101010 | 1100110011001011 |  |
| 1 | 1 | 1 | 0 | NOT | 1100110011001100 | 1010101010101010 | 0011001100110011 |  |
| 1 | 1 | 1 | 1 | CLEAR | 1100110011001100 | 1010101010101010 | 0000000000000000 | 0 |

Table-4 shows the Fourth 4 operations selected depending on select lines. The operation is performed on 2 input's A \& B and output is given out. The operations performed are SUBTRACTION, DECREMENT, NOT, and CLEAR.

## VI. CONCLUSION

This 16-bit reversible Alu is designed and implemented in Verilog using MODEL SIM ALTERA 6.6d. The main aim of the design in this paper is improve the ALU features by increasing it to 15 -operations and increase width to 16 -bit. For further research this ALU can be extended to 32 -bit and 64 -bit and more features can also be added.

This design is verified using Verilog which has constrains of input to output one-way functionality, if we can design the reversible logic circuit using tools which support 2-way functionality the reversible logic result can be simulated and analysed in much better.

## REFERENCES

[^0]
# International Journal of Advanced Research in Electrical, Electronics and Instrumentation Engineering 

(An ISO 3297: 2007 Certified Organization)

## Vol. 3, Issue 4, April 2014

[4] Feynman, R., "Quantum mechanical computers" Optics, News, and pp.: 11-20 11, 1985.
[5] Toffolli T, "Reversible computing", Technical Memo MIT/LCS/TM-151, MIT Laboratory for Computer Science (February), 1980.
[6] Fredkin, E. and T. Toffolli, "Conservative logic", Intl. J. Theoretical Phy., 21, pp.: 219-253, 1982.
[7] B.Raghukanth, B.Murali Krishna, M. Sridhar, V.G. SanthiSwaroop"A DISTINGUISH BETWEEN REVERSIBLE AND CONVENTIONAL LOGIC GATES", International Journal of Engineering Research and Applications (IJERA), Vol. 2, pp.148-151, Issue 2,Mar-Apr 2012.
[8] B. Raghu Kanth, B. Murali Krishna, "A COMPARITIVE STUDY OF REVERSIBLE LOGIC GATES", International Journal of VLSI \& Signal Processing Applications, Vol.2, Issue 1, PP: (51-55), Feb 2012.
[9]Vandana Shukla, O. P. Singh, G. R. Mishra, \& R. K. Tiwari, ASET, Amity University. "Design of a 4-bit 2's Complement Reversible Circuit for Arithmetic Logic Unit Applications", Special Issue of International Journal of Computer Applications ( 0975 - 8887), The International Conference on Communication, Computing and Information Technology (ICCCMIT), 2012.
[10] Md.Sazzad Hossain, Uttam Kumar Acharjee, \&NazmulHaque of MawlanaBhashani Science \& Technology University, "Design of Control Unit for Low Power ALU with a Barrel Shifter Using Reversible Logic", International Journal of Engineering, Business and Enterprise Applications (IJEBEA), 2013.
[11] Ann Susan Varghese, M Tech Student, Anjali V, assisstant Professor, Dept. of Electronics \& Communication, Mangalam College Of Engineering, Ettumanoor, Kottayam, India, "Design of 32 Bit Reversible ALU", , International Journal of Electronics \& Communication Engineering Research (IJECER), Page: 53-56, Vol. 1 Issue 3, August - 2013.
[12] Akanksha Dixit, VinodKapse , "Arithmetic \& Logic Unit (ALU) Design using Reversible Control Unit", International Journal of Engineering and Innovative Technology (IJEIT) Volume 1, Issue 6, Page: 55-60, June 2012

## BIOGRAPHY



Abdul Kareem Born in Mysore \& brought up in Ramanagaram, completed B.E in Electronics \& Communication from Ghousia college of engineering, Ramanagaram. Pursuing M -Tech in Electronics from HKBK College of engineering. Nagawara, Bangalore.


Mrs. E Kavitha obtained her B.E. and M.E. degrees from Bharathidasan University and Madras University in the year 1995 and 2000 with FCD. She is working as Professor PG-Coordinator in HKBK College of Engineering, Bangalore. She has 19 years of experience in Electronics and applied electronics. Her areas of interest are communication, wireless networks, optical networks and Image Processing.


[^0]:    [1] Y.syamala \& A.V.N Tilak , "Reversible Arthematic logic unit", IEEE conference, Page no: 207-211, 2009
    [2] R. Landauer, "Irreversibility and heat generation in the computing process", IBM J. Research and Development, 5(3), pp.: 183-191, 1961.
    [3] C. H. Bennett, "Logical reversibility of computation", IBM J. Research and Development, 17, pp.: 525-532, , November1973.

