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An Efficient SBST Scheme for Multiple-Memory based Multi-Processor Circuits

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ABSTRACT: To reduce the performance stemming caused due to the execution of branch instructions in pipelined architecture, BPU architectures play a very important role. Mainly BPU has got two mechanisms, BHT and BTB. In this paper a new mechanism which is a combined form of both is described and testing of this BPU memory is another constraint. The conventional method used to test the BPU memory was hardware testing scheme known as built-in self-test (BIST). In order to overcome the disadvantages of BIST software based testing scheme was developed, software based self-testing (SBST), where a separate entity is not required while testing a module. In this paper a new genetic algorithm is introduced as an efficient SBST scheme for testing BPU memory to make the complete system more efficient by reducing the redundancy of test patterns to the maximum extent by using various evolutionary strategies. A comparative study followed by a complete optimisation in terms of area, power and delay between the three architectures are done and in order to make it application specific the existing as well as the proposed systems are then implemented on a multiple memory based multiprocessor circuit. The very high speed IC hardware description language (VHDL) model of the complete system is simulated using the most efficient professional fault simulator ModelSim SE 6.2b.Also for the synthesis and analysis of HDL designs, for timing analysis and to obtain the design summary of various modules XILINX ISE 13.2 is used. Also the hardware implementation of the system is done using MATLAB and SPARTAN 6 FPGA kit.

KEYWORDS:testing, multiprocessor, built-in self-test, software based self-test, genetic algorithm, evolutionary strategies, branch history table, branch target buffer, branch prediction unit, very high speed IC hardware description language.

I.INTRODUCTION

Nowadays pipelined architectures/processors play a very important role due to the significant improvement in their performance. Pipelining has got mainly four stages: - fetch, decode, execute and write back. When a data is being fetched after completing the decoding stage that is before entering to the execution stage the processor is capable of fetching a new instruction. But after reaching the second or third stage only the processor will do a prediction about whether the data was previously executed or not. If it results in a wrong prediction, an increased level of performance stemming will occur. This negative effects stemming from branch instructions can be avoided by branch prediction unit (BPU). The BPU has got mainly two architectures, Branch history table (BHT) and branch target buffer (BTB). The BHT is just like a look up table (LUT) which stores the results of previously executed conditional branches. Since the redundancies of test patterns are very high and it offers only the testing of very few test patterns architecture came known as Branch target buffer (BTB). BTB mainly forecasts about the target address. Since the computational density is very high it's very necessary to adopt some other method for BPU design. Here in this paper a combined technique is described which includes all the advantages of BHT as well as BTB and also averages the drawbacks of both. The testing of this BPU memory is another constraint and here a new genetic algorithm which is an efficient replacement for both software based self-test (SBST) as well as built-in self-test (BIST) is introduced. Very frankly, the new test schemes do not necessarily aim to substitute other testing schemes (scan chains, BIST) but rather to supplement the



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same by adding more test quality at a low cost. To make it application specific the whole system is implemented using a multiple memory based multiprocessor circuit. The main key problem regarding pipelining is branches.

Now consider a code sequences,

If (p>10) { q = r; } else { q = s; } This compiles into something like Cmp p, 10; p>10? b1e L1 movr,q; q=r br L2 L1: movs,q; q=s L2:----

By the time the conditional branch at line 2 reaches the execute stage in the pipeline, the processor must have already fetched and decoded the next couple of instructions. But the processor may be ambiguous about which instruction should it fetch and decode, either to fetch and decode the if branch (lines 3, 4) or else branch (line 5)? It won't really know until the conditional branch gets to the execute stage, but in the case of deeply pipelined processor that might be several cycles away and it can't afford to just wait. For every six instructions on average the processor encounters a branch, and if it was to wait several cycles at every branch then most of the performance gained by using pipelining in the first phase would be lost. Here is the importance for a BPU. The branch predictor attempts to avoid the waste of time by trying to guess whether the conditional jump is most likely to be taken/not taken

II. RELATED WORKS

In [2], an O/S-enabled framework Deamon Guard for on demands (selective) SBST to support fault recovery capabilities. Towards this goal, the paper proposes an efficient check pointing and rollback recovery mechanism which, upon fault detection, can restore the system to the most recently valid Correct state and resume the normal operation assuming disabling of the faulty Core, thereby leading to a healthy (but degraded) system. In [3] the basic testing method to verify a failure free status automatically, without the need for automatically applied test stimuli (other than power and clock), and without the need for the logic to be part of running system. Since the overall cost associated with the circuit implemented using BIST is much less in comparison to all other testing techniques and it is used as the most common testing scheme in earlier days. The SBST an SBST program development methodology is proposed for online testing of small cache memories in microprocessors, also the methodology have applied a March tests to detect storage faults and comparison faults when the cache comprises content addressable memories (CAM) arrays [4]. Another approach describes a key idea of SBST is to exploit on-chip programmable resources to run normal programs that test the processor itself. The processor generates and applies functional-test patterns using its native instruction set, virtually eliminating the need for additional test- specific hardware [5].the [6] is all about a new method to generate a proper test program to be executed by a processor in order to check whether the BPUs implementing the Branch History Table (BHT) architecture works correctly, based on a purely functional approach, i.e., it does not require any knowledge about the actual implementation of the circuitry it is intended to test, nor on the adopted semiconductor technology. A multithreaded (MT) SBST methodology generates an efficient multithreaded version of the test program and schedules the resulting test threads into the hardware threads of the processor to reduce the overall test execution time and on the same time to increase the overall fault coverage [7]. In [8] after the MT SBST a new testing scheme automates the test by employing the deterministic behaviour of the BPU based on a predetermined sequence of branches, a fault detection implementation using a special DFT logic, and an alternative solution utilizing a low cost Multiple Input Signature Register (MISR). The multiprocessor sharing multiple memory and also using the common test vector for all the sub modules to reduce the redundancy and rolling back the system to its maximum used check point [9].In [10] the online testing of chip multiprocessor can be done on both basis hardware/software mechanisms. The software architectures are used for the detection of online error detection and the recovery of the same can be done by the use of multicore processors [11].



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III.SYSTEM MODEL

The existing system is all about a branch prediction module (BHT, BTB) design. And the testing of BPU memory which is embedded inside a processor using software based testing scheme SBST, and also the conventional testing scheme which is hardware based known as built-in self-test (BIST).

3.1. Branch History Table

A BHT is based on a data structure which stores the result (Taken or Not Taken) of previously executed conditional branches. The goal of a BHT is simply to forecast whether a given conditional branch result will be taken or not. The BHT is accessed during the Decode stage each time the instruction processed by the stage is a conditional branch.



Fig.1BHT Module

This BHT is designed using a Look-Up Table LUT) that is nothing but an 8 bit ROM of size 256, an 8 bit state machine which acts like a decision block with input equals the output from ROM as shown in figure 1. The decision block checks the LSB of input and accordingly it releases the output value, checking the first 2 MSB bits of this output it is possible to predict whether the test pattern are taken or not taken. If the 2 bit MSB is "10" it is predicted as not taken and "01" means taken. The rest of the bits of the output resemble the LSB values of input. The output from LUT which is fed as input to the decision block is nothing but the output from 8 bit LFSR which is equal to the value which is stored in the RAM for the corresponding address. In order to check the redundancy of test pattern it's necessary to compare the output from the LUT as well as the output from decision block using an XOR gate, which is fed back to the LUT. Since it checks only few test patterns and also creates redundancy BHT based BPU mechanism came.

3.2. Branch Target Buffer

The BTB adopts a more aggressive approach to branch prediction, just by checking the target address it will come to know whether the data has been previously executed or not. Moreover the instructions are accessed during the fetch stage itself. For this design an 8-bit linear feedback shift registers (LFSR), An 8-bit RAM of size 256, where the output from this LFSR is fed as input address to RAM. The output from the 8 bit LFSR is fed as one of the input to the XOR module which act as the present value; also the other input to this XOR module is an array of values in order to check for previous history while doing prediction. The output from this XOR module is fed as input address to the RAM memory. Finally, by checking this address value it is easy to predict whether the branch has been previously taken or not. If it is zero value, it is taken and otherwise it will give the present value. The figure 2 shows the block diagram of branch target buffer (BTB). The computational density is very high for BTB even though it is less redundant.



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Fig.2BTB Module

3.3 BIST

Built-in self-test is one of the conventional testing scheme. If BIST is used to test a module it is necessary to add a separate entity while doing the test since it is purely hardware based. The figure 3 shows built-in self-test (BIST) block diagram. The de-bouncer actually consists of 6 inputs and outputs respectively. When a switch is being pressed and hold during that particular time it reads many value, so in order to take a single value as the output it is used. By giving the one of the switch output value of de-bouncer and a reset value as inputs the teller produces an output which in turn fed as one of the input to multiplexer, the output is nothing but from the LFSR output. Using one of the switch output as select line one of the input value will get switched



Fig.3BIST Module

to the output. The complete lower blocks are for the generation of seed value for the BIST. The output from de multiplexer is fed as input to the BIST. The BIST consists of mainly two blocks: -LFSR, SIRS. From this block the error value after testing is received. It has got many disadvantages like high area overhead, increased power consumption also due to large signal routing paths.

3.4. SBST

The Software Based Self-Test (SBST) is actually a software tester code which is embedded inside the processor for testing the BPU memory more efficiently. Compared to conventional testing scheme extra hardware is not needed



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while testing a module. Actually a small separate code can be used for SBST implementation and it is then embedded into the main program module for testing the memory. Normally the memory testing using SBST consists of two process states. By giving clock, idle=0 it acts like a normal RAM, during the second process states by giving second clock tclk, idle=1 it starts testing the memory and give the error value as well as non-redundant test patterns as the output.

IV.PROPOSED SYSTEM

4.1. Combined architecture

Since the single architectures of BPU have got their own disadvantages, it is essential to find out a new mechanism which combines all the advantages of BHT, BTB as well averages the demerits of both. The redundancy of final test patterns is negligible when compared to single architecture. Also testing of all the test patterns are carried out consuming very less time with maximum efficiency.



Fig. 4 combined architecture

In figure 4, the block diagram of combined architecture of BPU is shown. It consists of both single architecture modules BHT, BTB. The outputs obtained from both are sorted to obtain non-redundant test patterns and fed as data input to the RAM_SBST module which in turn used for the testing of the memory.

4.1.1. Comparative study of BPU mechanisms

The combined module is actually a combination of advantages of both BTB, BHT. Even though the time consumption of all the three mechanisms are comparable it consumes only very few power compared to other two and area compared to BHT is Also very less. The Table 1 shows the comparative study of all the three BPU mechanisms in terms of area, power and delay.

BPU Mechanism	Area (No of slices)	Delay (ns)	Power (mw)
BHT	99	7.141	204
BTB	63	7.141	204
Combined Module	64	7.141	80

Table 1Study of BPU mechanisms

4.2. Genetic Algorithm

The algorithm described in this paper is a modified one for the existing SBST. This genetic algorithm is all about creating higher populations using various evolutionary strategies, thus reducing the redundancy of test patterns to the maximum extent. Here the evolutionary strategies are nothing but the patterns coming out from an 8-bit linear feedback shift registers (LFSR).keeping these values as reference the proximity of each value stored in an LUT is checked



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continuously. And according to the proximity each test patterns from the main LUT is subdivided and stored into small eight groups each having the same proximity. In order to avoid the redundancy further and also to make the SBST testing scheme more efficient, each test patterns from the main LUT is checked with those test patterns which are stored inside eight LUTs and if it shows any resemblance one of the value from the respective group is taken as output value only for once.

4.3. Comparative Study

genetic Project Status (05/04/2016 - 13:12:17)										
Project File:	geneticc.xise	Parser Erro	rs:	No Errors						
Module Name:	genetic	Implement	ation State:		Programming File Generated					
Target Device:	xc3s250e-4vq100	• Erro	rs:		No Errors					
Product Version:	ISE 13.2	• War	nings:		1/1 Warning	as (0 new)				
Design Goal:	Power Optimization	• Rou	ting Results:		All Signals C	ompletely R	outed			
Design Strategy:	strategy1	• Tim	ing Constraints		All Constrain	nta Met				
Environment:	System Settings	• Fina	I Timing Score:		0 (Timing R	eport)				
	Device Uti	lization Sun	mary							
Logic Utilization		Used	Available	Utili	zation	Note(s)				
Number of Slice Elip Elops	an l		4,096		1%					
Number of 1 input LUTs		61	1,896		1%					
Number of occupied Sirce	9	19	2,118		27%					
Number of Slices conta	ining only related logic	19	19		100%					
Number of Slices conta	ining unrelated logic	0	10		0%					
Total Number of 4 Input L	UTS	70	4,895		1%					
Number used as logic		61								
Number used as a rout	re-thru	9								

Fig.5 Design Summary of Genetic Algorithm

After doing the comparative study with the conventional testing scheme (BIST), existing method (SBST), it is more efficient to use genetic algorithm to test the memory inside the processors. Complete optimisation is obtained in terms of area, power and delay. The area utilisation in terms of number of LUTs, slices are very less for genetic algorithm compared to other two schemes shown in figure 5

GENERATED AFTER PLACE-and-ROU	TE.		
Clock Information:			
Clock Signal	Clock buller(FF name)	Load	11
clk	BUFGP	38	11
Asynchronous Control Signals Inform 	willon: md in this design h Frequency: 163.900MHz) c clock: 5.781no er clock: 5.781no er clock: 4.310ns :: No path found		

Fig.6Timing Report of Genetic Algorithm

Also the power consumption is much less for this genetic algorithm as shown in figure 7 and very high for the BIST. The delay factor is also comparatively low for this proposed algorithm as shown in figure 6, and on a balanced level for the existing scheme (SBST).



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Fig.7 Power Analysis of Genetic Algorithm

The table 2 shows the comparative study of all the three testing mechanisms in terms of area, power and delay. It is clear that the area as well as power consumption is comparatively very less for the proposed system and also the delay is less even though it doesn't make much difference when compared to others.

Testing schemes	Area	Delay	Power
	(No of slices)	(ns)	(mw)
BIST	104	8.015	205
SBST	64	7.141	80
Genetic Algorithm	50	6.098	51

Table 2	Study	of Testing	Schemes
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4.4. Application

The multiple memory based multi-processor platform can be used for the implementation of existing as well as proposed system. Here a simple ALU can perform the role of this multiprocessor by designing two ALU's, two Ram's. A single ALU can access both the memories at a time but the other ALU can only access the memories after checking the select lines x1, x2, x3, x4 as well as the value of the variable BZY1, BZY2 as shown in figure 8.



Fig.8 Multiple memories based multiprocessor circuit



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The existing system, SBST can be implemented by taking the first sequence from final output which is then fed as one of the input to the ALU and taking the second sequence which in turn acts as the second input of ALU. Since the select line of ALU is half the number of bits of SBST output test patterns, only the 4 bits from LSB of the third sequence is fed as select line to the ALU.At the same time the proposed system, genetic algorithm can also be implemented using multiple memory based multiprocessor. One of the outputs from the genetic algorithm is fed as one of the input to the ALU. Since the ALU consists of two inputs the other one is kept as a constant, and here it is five. And the outputs are stored accordingly to the output variable of ALU.

V.SIMULATION RESULTS

The hardware is modelled using VHDL coding. Simulation results are verified in Model Sim 6.2b. and it is synthesized and implemented in Xilinx ISE 13.2, using SPARTAN 3E. The output waveforms of BHT is shown in figure 9 in which the outputs y3 from LUT and outtemp2 from decision block are fed as input to XOR and the final output outtemp3 gives the non-redundant test patterns.



Fig.9 Output waveform of BHT

The figure 10 shows the output waveform of BTB y, y1 acts as the two inputs of XOR module and the output from the same is y2 which is fed as the input address to the RAM module. According to the value of we the output is decided, whether to switch the output to d_o or to y.

👌 bitre_anck	1								
👌 /bthrr_api/st	1								
👌 /bitnw_apr/xe	1								
idoni_apid_									
c_buq_mddl 🙌	1001111	OICIID							
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0.4	979)								

Fig.10 Output waveform of BTB

The figure 11 shows the output waveform of BIST. The lfsrdatain is the input seedvalue generated from MUXOUT which act like a demultiplexer is fed to the BIST module. Signature out, lfsrout, ledbist are the outputs respectively from BIST. Also showing the outputs as well as inputs of functional output block, fout.



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Fig.11output waveform of BIST

The output waveform of combined architecture is shown in figure 12. The two outputs outtemp,out1 gives the maximum non redundant testpatterns compared to single architecture as well as error value after testing the memory. The value of d_i is the XOR'ed output of both BHT, BTB which is then fed as input to the main module.

4	/combine_api/g2/clk	1										
\$	/combine_api/g2/tclk											
4	/combine_apt/g2/we											
4	/combine_apr/g2/ad	0000001								00000		
\$	/combine_apt/g2/ideal											
I\$	/combine_apr/g2/ou		0000	101	10100100				01010010			
4	/combine_aps/g2/out1											
₩	/combine_apr/g2/d		UUUU	UUUU								
₩	/combine_apr/g2/d_i		10001	00								
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I\$	/combine_apr/g2/x		10010	101			00101011				01010110	
I\$	/combine_apr/g2/te		10100	100	01010010				10101001			
\$	/combine_api/g2/11											
\$	/combine_api/g2/12											
ł	/combine_apr/g2/y2		C1010	010			10101001				<u>11010100</u>	
\$	/combine_apı/g2/co		(
			_									

Fig.12 Output waveform of combined architecture

The figure 13 above shows the output waveform of genetic algorithm.lut0 to lut7 shows the 8 subgroups of main lut(lutstd).The dataout shows the selection of subgroup done with reference to the main lut which is grouped according to te proximity between lutstd,lfsr.

4	/genetir_ap/dk]					i i	i ii	i i	İ I	1 11	
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*		101001010 001100	(01001010.0011)	to11 001100°O	00110001.0011	ໝັກມານມ	JUU CO101110	00101101 001	0 100 0010101	1 00101010000	i da con con can	00 001 <mark>0</mark> 011
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Fig.13 Output waveform of Genetic

The above figure 14 shows the implementation of SBST using multiprocessors.p1,p2,q1,q2,r1,r2 are inputs as well as select lines of ALU which are actually the output test patterns of SBST.O1,O2 are the final outputs of ALU respectively after implementing SBST.



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Fig.14 Implementation of SBST

The implementation of genetic algorithm is shown in figure 15 below. The output of genetic algorithm is fed from testt to ALU, one of the input of ALU and the other one is noting but a constant value "00000101.cout is the final output.



Fig.15 Implementation of Genetic Algorithm

VI. HARDWARE IMPLEMENTATION

The hardware implementation of the proposed system that is genetic algorithm can be done using the software MATLAB 9.1, SPARTAN 6 FPGA kit. After checking the similarity of test patterns between those stored inside the main LUT and eight groups, each test patterns is taken out respectively from the groups to minimize the redundancy as well as to reduce the complexity while testing the modules. So here according to the selection of groups from lut0 to lut7 the LED's will toggle as shown in figure 16.The same way it is possible to implement existing as well as conventional schemes using FPGA. The SBST can be implemented by showing whether there is error or not after testing the memory by toggling the LED. Likewise BIST can also be implemented after giving the pin values of switch inputs and clock correctly and the toggling of LED's shows the output test patterns and also the error value.



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Fig.16 Hardware implementation

VII.CONCLUSION

VLSI design of an efficient BPU mechanism is proposed for hybrid or superscalar architectures. Designed the combined architecture of BPU which includes both BHT, BTB using VHDL coding. Thus the performance stemming can be reduced to maximum in a super scalar architecture eliminating the drawbacks of single architecture of BPU. Also a comparative study between all the three BPU mechanisms is done. And the main focus is given to the new genetic algorithm and the comparative study with other two testing schemes (BIST, SBST) in terms of area, power, and delay. By using this algorithm, it is possible to cover maximum number of test patterns with negligible redundancy and with very less power as well as area. The existing systems as well as the proposed systems are also implemented using multiple-memory based multi-processor. Simulated in Model SIM 6.2b and synthesized using Xilinx ISE 13.2.The hardware implementation can be done using the software MATLAB 9.1 and SPARTAN 6 FPGA kit.

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